

**Unit: III**  
**Lecture: 4**  
**Relational Algebra**  
**(Part-II)**

**Relational Oriented Operation**

➤ **Selection or Restriction Operation:** The selection operation is a unary operation. This is used to find the tuples of relations.

- **It is denoted by sigma ( $\sigma$ ).**
- e.g. if we want all the employees having salary more than 9000 from relation employee, the query is:

$\sigma$  salary >9000(Employee)

**Result:**

EID	Name	Salary
1E	John	10000
5E	Nile	15000

➤ **Projection Operation:** The Projection operation is a unary operation which applies only on a single relation at a time. Project operation is used to select columns of table.

- **It is denoted by pi ( $\pi$ )**
- e.g. if we want all the names of employees and their salary from relation Employee, then query is:

$\pi$  name,salary (Employee)

**Result:**

Name	Salary
John	10000
Ramesh	5000
Smith	8000
Jack	6000
Nile	15000

➤ **Natural Join Operation:** Natural join ( $\bowtie$ ) is a binary operation that is written as ( $r \bowtie s$ ) where r and s are relations. The result of the natural join is the set of all combinations of tuples in r and s that are equal on their common attribute names.

➤ **Left Outer Join( $R \ltimes S$ ):** All the tuples from the Left relation, R, are included in the resulting relation. If there are tuples in R without any matching tuple in the Right relation S, then the S-attributes of the resulting relation are made NULL.

➤ **Right Outer Join: ( $R \ltimes S$ ):** All the tuples from the Right relation, S, are included in the resulting relation. If there are tuples in S without any matching tuple in R, then the R-attributes of resulting relation are made NULL.

➤ **Full Outer Join: ( $R \ltimes S$ ):** All the tuples from both participating relations are included in the resulting relation. If there are no matching tuples for both relations, their respective unmatched attributes are made NULL.